

Introduction

Joint Degree Matrix

- Joint Degree Matrix (JDM) for a graph $G(V, E)$:

$$JDM(k, l) = \sum_{u \in V_k} \sum_{v \in V_l} 1_{\{(u,v) \in E\}}$$

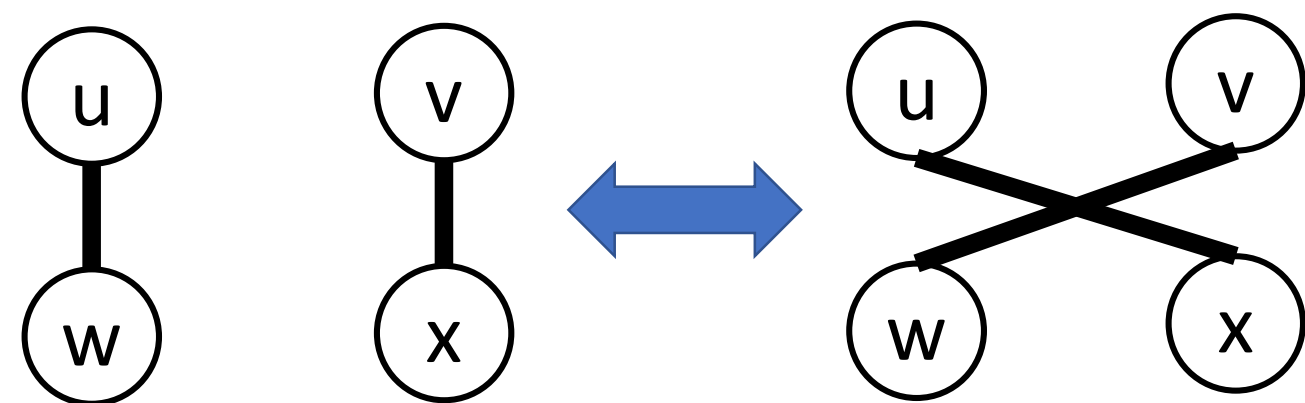
$$V_i := \{v \in V \mid \text{degree}(v) = i\}$$

JDM Realizations

- Simple graphs with the same target Joint Degree Matrix (JDM) are called realizations of that JDM

JDM-preserving double-edge swap

- Given four distinct nodes (u, v, x, w) and $(u,w), (v,x)$ edges [and $(u,x), (v,w)$ are not edges]. Deleting $(u,w), (v,x)$ and adding $(u,x), (v,w)$ preserves JDM if $u, v \in V_i$

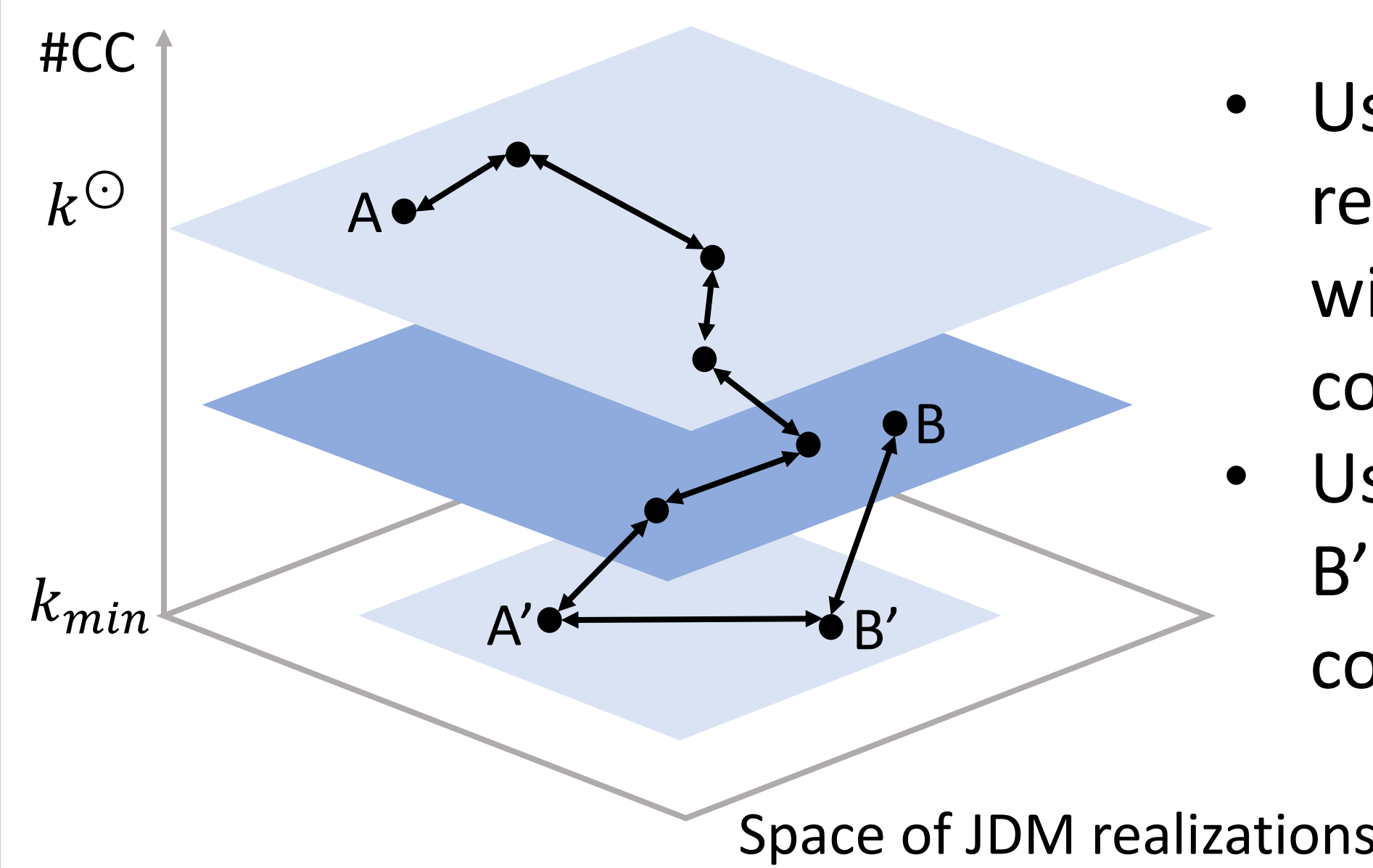


Prior Work

- Space of realizations is connected over JDM-preserving double-edge swap operation:
 - Symmetric difference based proof [Amanatidis '15]
 - Balanced realization based proof [Czabarka '15]
- Construct JDM realization with single connected component [Amanatidis '15, '18]

Main result

- Theorem.** The space of simple, undirected graphs with a target JDM and no more than k^\odot number of connected components is connected under a sequence of JDM-preserving double-edge swaps.

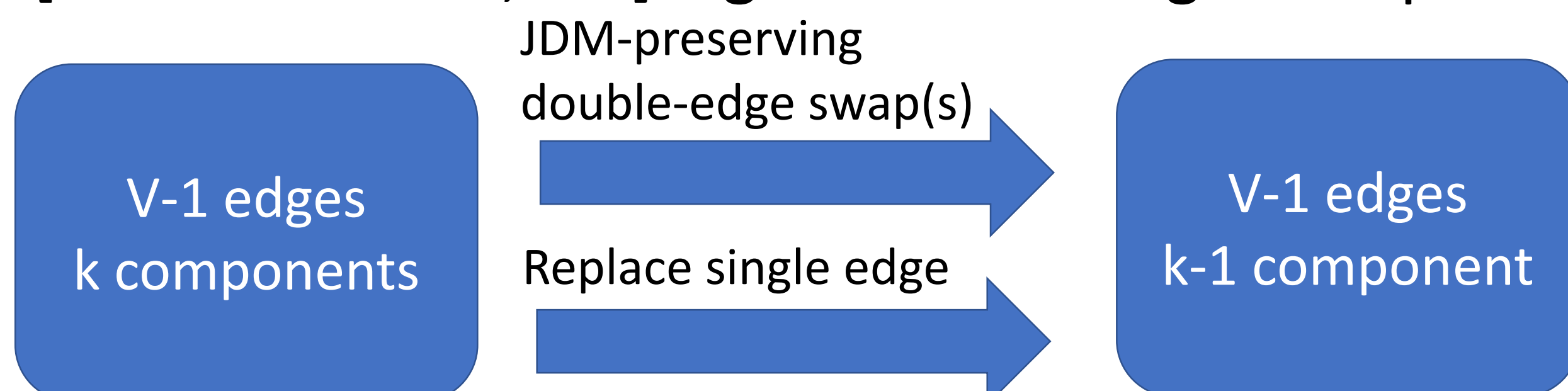


- Use Lemma 1 to transform JDM realizations A and B to A' and B' with minimum number of connected components
- Use Lemma 2 to transform A' to B' while preserving number of connected components.

- Lemma 1.** For a realizable JDM, there exists a JDM-preserving double-edge swap sequence that transforms any realization to a realization with minimum number of connected components, such that there is no double-edge swap that increases the number of connected components.
- Lemma 2.** The space of JDM realizations with minimum number of connected components, k_{min} , is connected under JDM-preserving double-edge swaps.

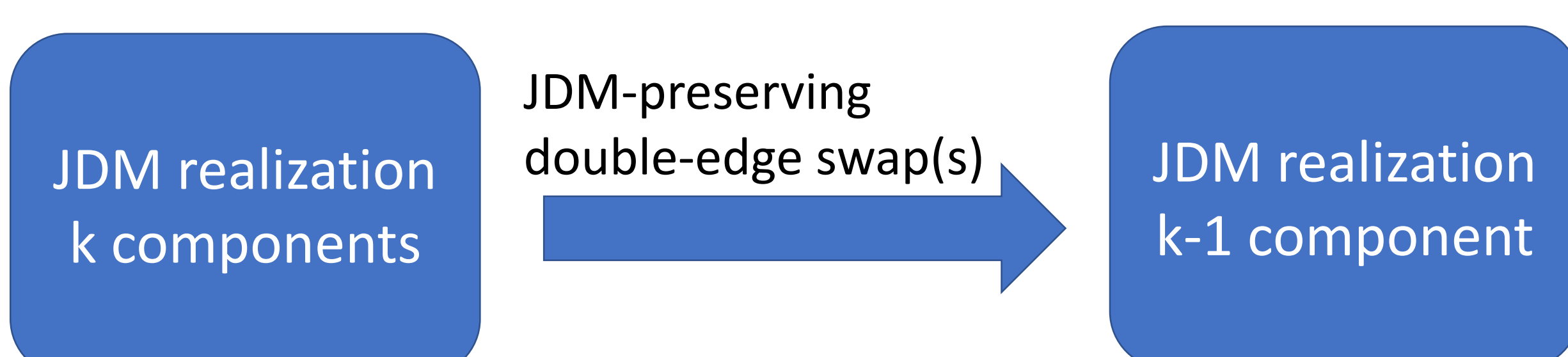
Proof Sketch of Lemma 1

- [Amanatidis '15, '18] algorithm for single component:

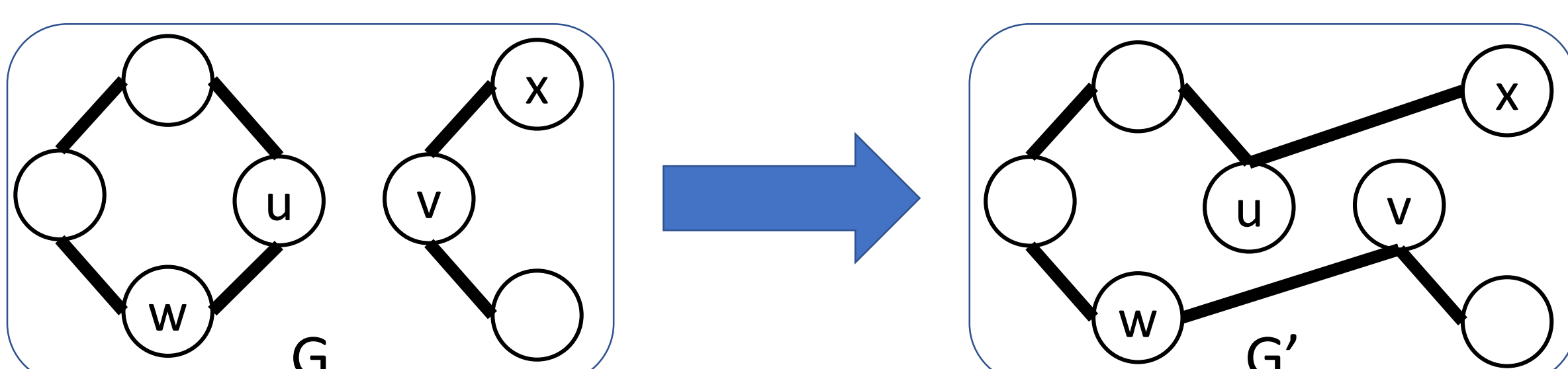


Repeat until either $k=1$, a spanning tree for target JDM or produce certificate that single connected component realization doesn't exist

- Our algorithm for minimum number of components:



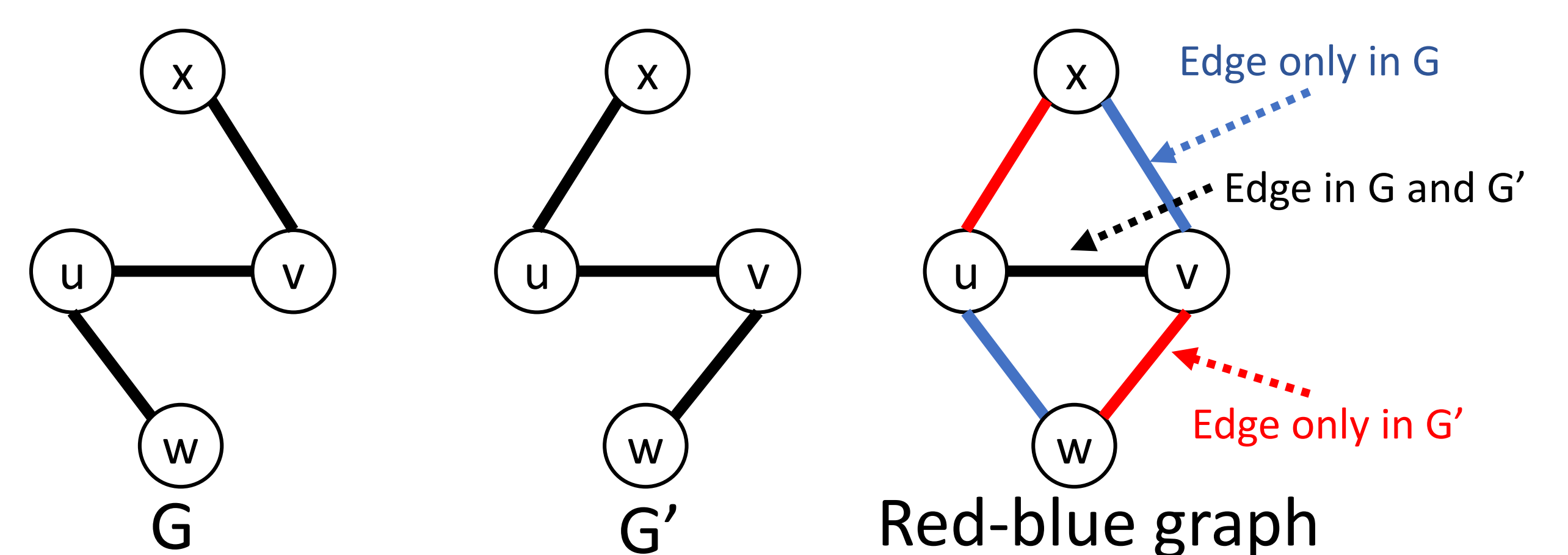
Repeat until either $k=1$ or produce certificate that $k-1$ components is not realizable thus current realization has minimum number of components



Key idea: JDM-preserving double-edge swap using $(u,w), (v,x)$ breaks a cycle to merge two connected components ($u, v \in V_i$)

Proof sketch of Lemma 2

- Symmetric difference of G and G' (red-blue graph):
 - Exists a red-blue alternating cycle decomposition



- Goal: Given two realizations G, G' , each with k_{min} components, find JDM-preserving double-edge swap(s) to reduce the symmetric difference without increasing number of connected components in G or G'

- Example case:**

